

CS 380 - GPU and GPGPU Programming

Lecture 24: CUDA Memory, Pt. 3

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Reading Assignment #13 (until Nov 27)



Read (required):

- Optimizing Parallel Reduction in CUDA, Mark Harris,
<https://developer.download.nvidia.com/assets/cuda/files/reduction.pdf>
- Programming Massively Parallel Processors book, 4th edition
Chapter 8: Prefix Sum (Scan) – work efficiency in parallel algorithms

Read (optional):

- Programming Massively Parallel Processors book, 4th edition
Chapter 7: Reduction
- GPU Gems 3 book, Chapter 39: Parallel Prefix Sum (Scan) with CUDA
https://developer.nvidia.com/gpugems/GPUGems3/gpugems3_ch39.html
- Faster Parallel Reductions on Kepler, Justin Luitjens
<https://devblogs.nvidia.com/parallelforall/faster-parallel-reductions-kepler/>

CUDA Memory: Shared Memory

Memory and Cache Types



Global memory

- [Device] **L2 cache**
- [SM] **L1 cache** (shared mem carved out; *or* L1 shared with tex cache)
- [SM/TPC] **Texture cache** (separate, or shared with L1 cache)
- [SM] **Read-only data cache** (storage might be same as tex cache)

Shared memory

- [SM] Shareable only between threads in same thread block
(Hopper/CC 9.x: also thread block clusters)

Constant memory: Constant (uniform) cache

Unified memory programming: Device/host memory sharing

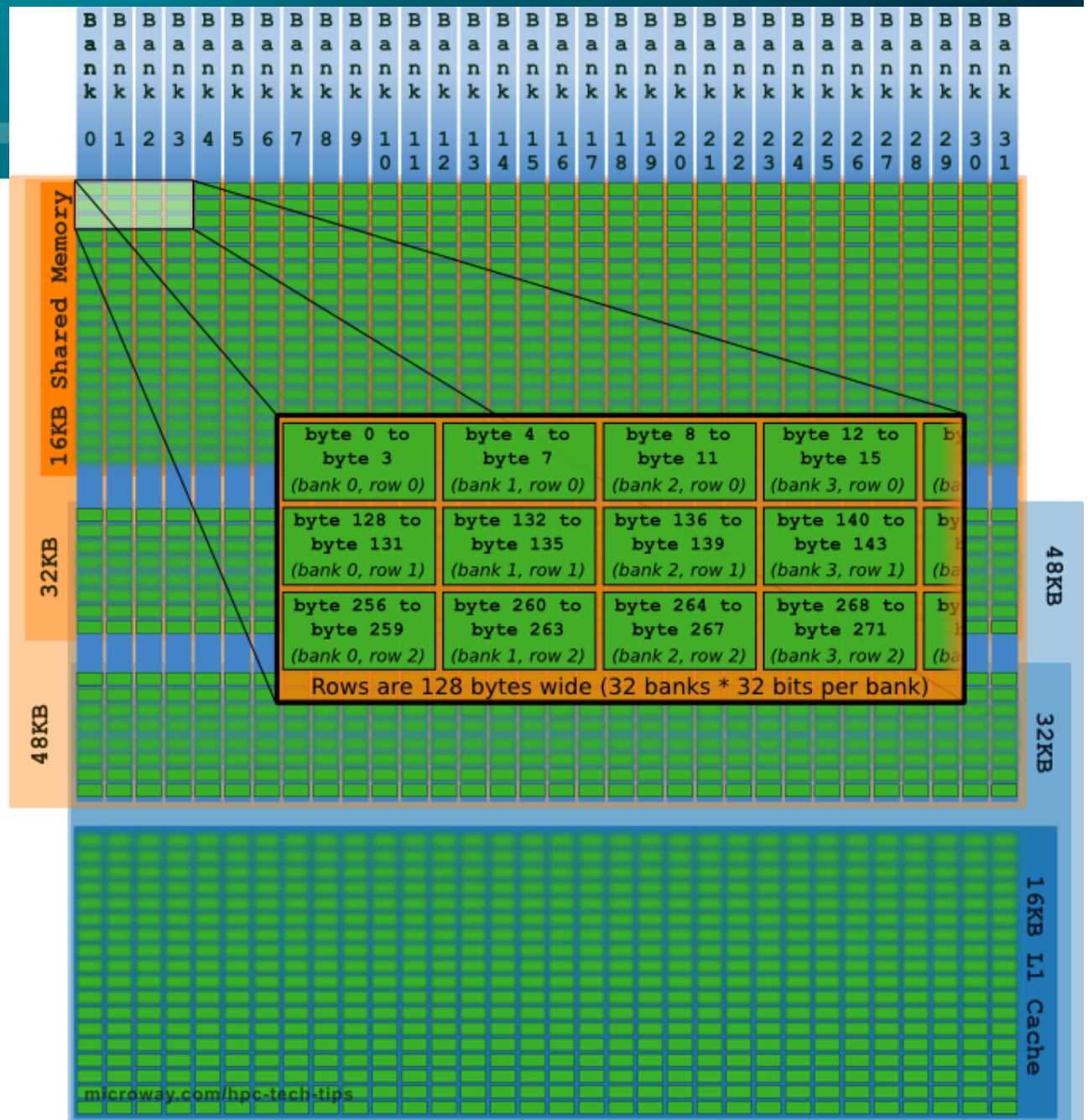
Memory Banks

Fermi/Kepler/Maxwell
and newer:

32 banks

default:
4B / bank

Kepler or newer:
configurable
to 8B / bank



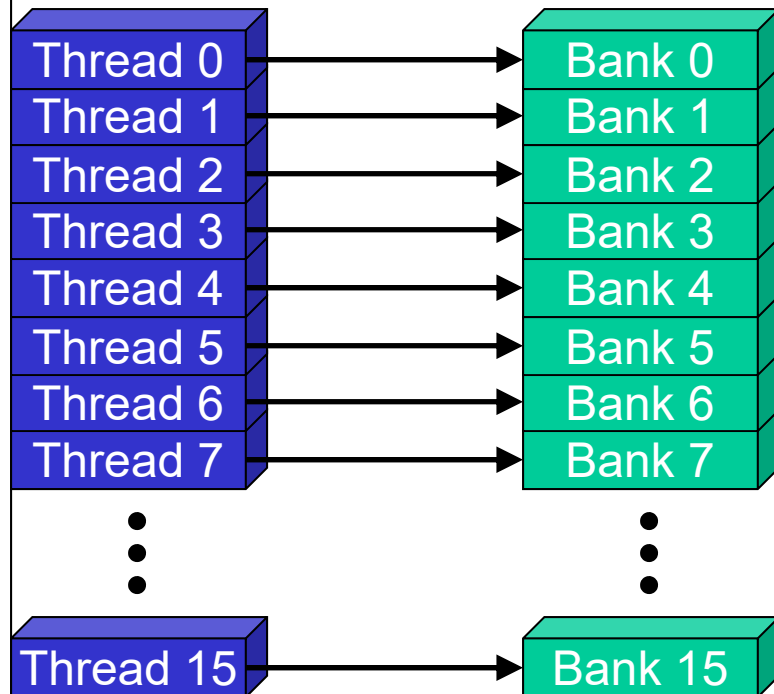
Shared Memory

- **Uses:**
 - Inter-thread communication within a block
 - Cache data to reduce redundant global memory accesses
 - Use it to improve global memory access patterns
- **Performance:**
 - smem accesses are issued per warp
 - Throughput is 4 (or 8) bytes per bank per clock per multiprocessor
 - **serialization:** if N threads of 32 access different words in the same bank, N accesses are executed serially
 - **multicast:** N threads access the same word in one fetch
 - Could be different bytes within the same word

Bank Addressing Examples

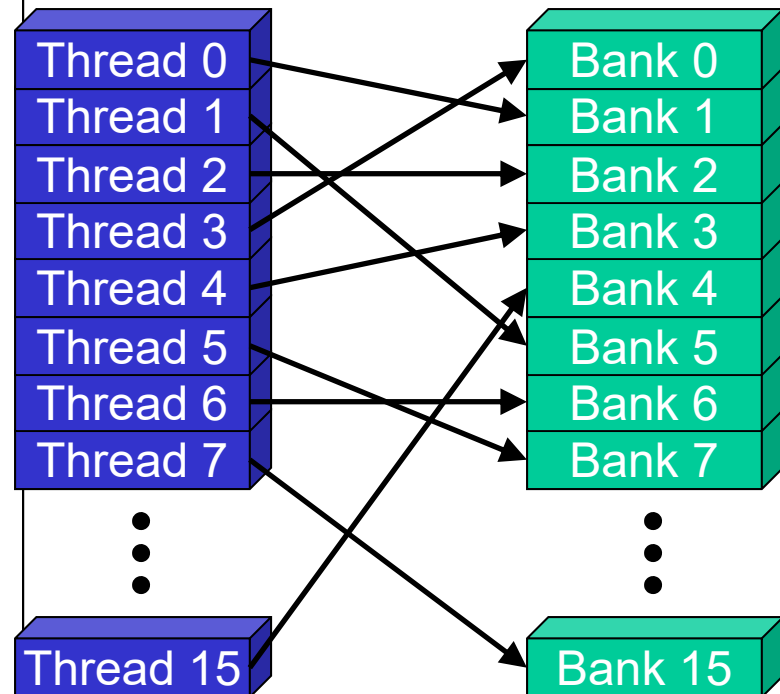
- No Bank Conflicts

- Linear addressing
stride == 1



- No Bank Conflicts

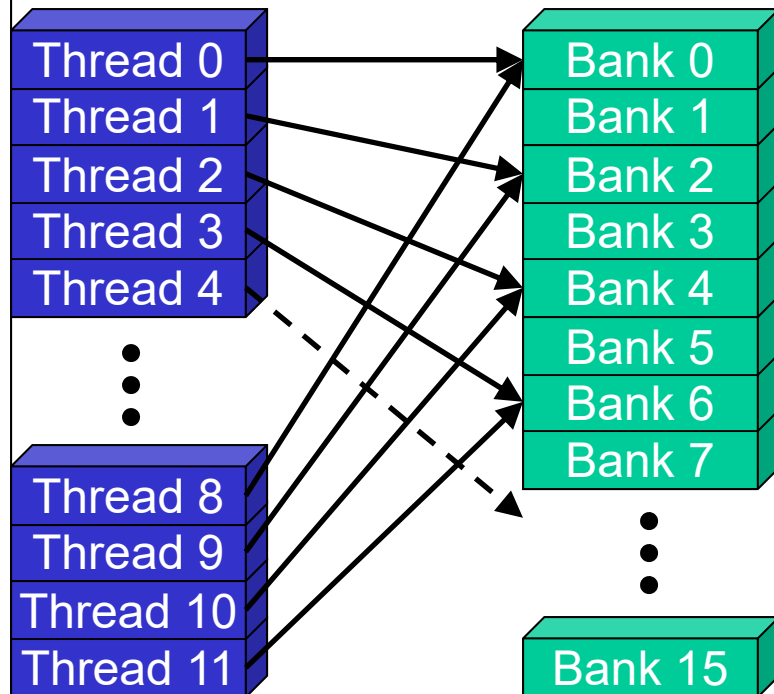
- Random 1:1 Permutation



Bank Addressing Examples

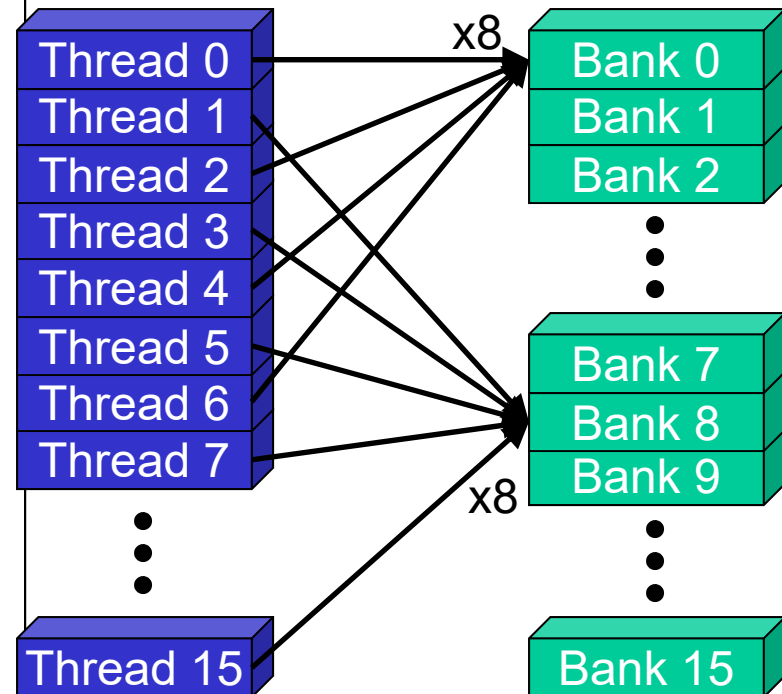
- 2-way Bank Conflicts

- Linear addressing
stride == 2



- 8-way Bank Conflicts

- Linear addressing
stride == 8



How addresses map to banks on G80

- Each bank has a bandwidth of 32 bits per clock cycle
- Successive 32-bit words are assigned to successive banks
- G80 has 16 banks
 - So bank = address % 16
 - Same as the size of a half-warp
 - No bank conflicts between different half-warps, only within a single half-warp

**Fermi and newer have 32 banks,
considers full warps instead of half warps!**

Shared Memory Bank Conflicts

- **Shared memory is as fast as registers if there are no bank conflicts**
- **The fast case:**
 - If all threads of a half-warp access different banks, there is no bank conflict
 - If all threads of a half-warp access the identical address, there is no bank conflict (broadcast)
- **The slow case:**
 - Bank Conflict: multiple threads in the same half-warp access the same bank
 - Must serialize the accesses
 - Cost = max # of simultaneous accesses to a single bank

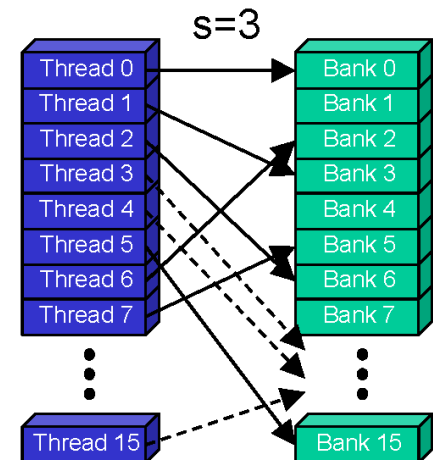
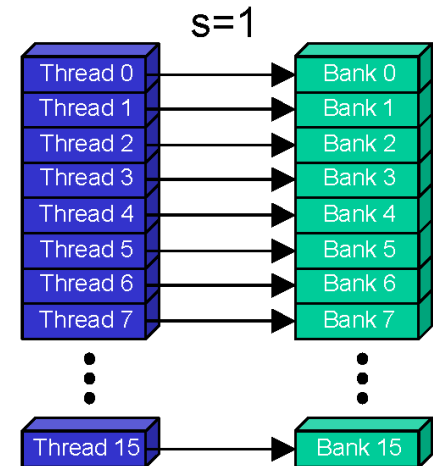
full warps instead of half warps on Fermi and newer!

Linear Addressing

- **Given:**

```
__shared__ float shared[256];  
float foo =  
    shared[baseIndex + s * threadIdx.x];
```

- **This is only bank-conflict-free if s shares no common factors with the number of banks**
 - 16 on G80, so s must be **odd**



Data Types and Bank Conflicts

- This has no conflicts if type of shared is 32-bits:

```
foo = shared[baseIndex + threadIdx.x]
```

- But not if the data type is smaller

- 4-way bank conflicts:

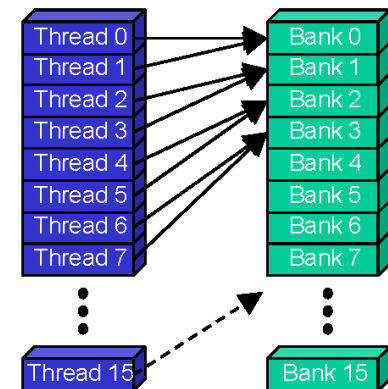
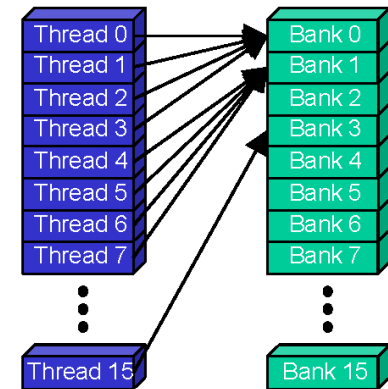
```
__shared__ char shared[];  
foo = shared[baseIndex + threadIdx.x];
```

but no problem on Fermi or newer: multi-cast!

- 2-way bank conflicts:

```
__shared__ short shared[];  
foo = shared[baseIndex + threadIdx.x];
```

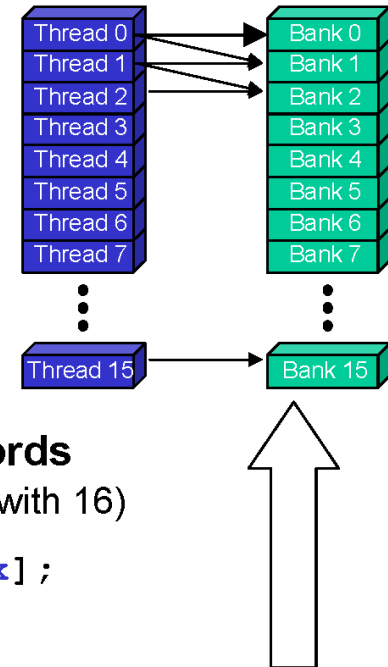
but no problem on Fermi or newer: multi-cast!



Structs and Bank Conflicts

- **Struct assignments compile into as many memory accesses as there are struct members:**

```
struct vector { float x, y, z; };  
struct myType {  
    float f;  
    int c;  
};  
__shared__ struct vector vectors[64];  
__shared__ struct myType myTypes[64];
```



- **This has no bank conflicts for vector; struct size is 3 words**
 - 3 accesses per thread, contiguous banks (no common factor with 16)

```
struct vector v = vectors[baseIndex + threadIdx.x];
```
- **This has 2-way bank conflicts for myType;**
(each bank will be accessed by 2 threads simultaneously)

```
struct myType m = myTypes[baseIndex + threadIdx.x];
```

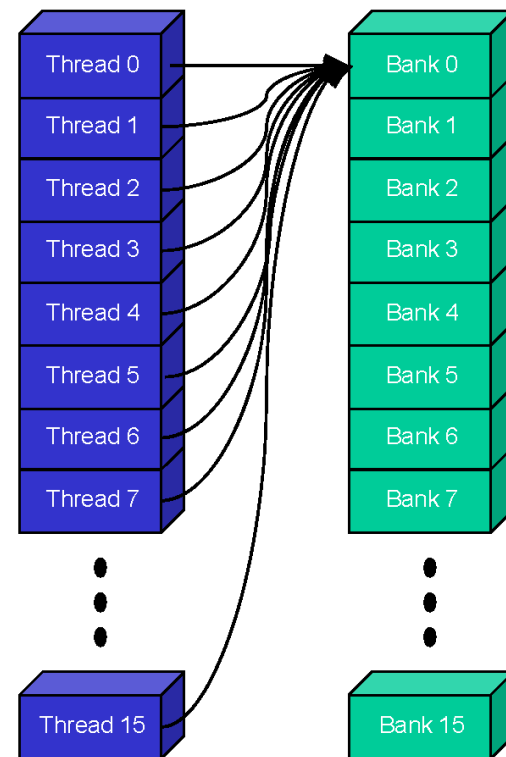
Broadcast on Shared Memory

- Each thread loads the same element – no bank conflict

```
x = shared[0];
```

- Will be resolved implicitly

multi-cast on Fermi and newer!



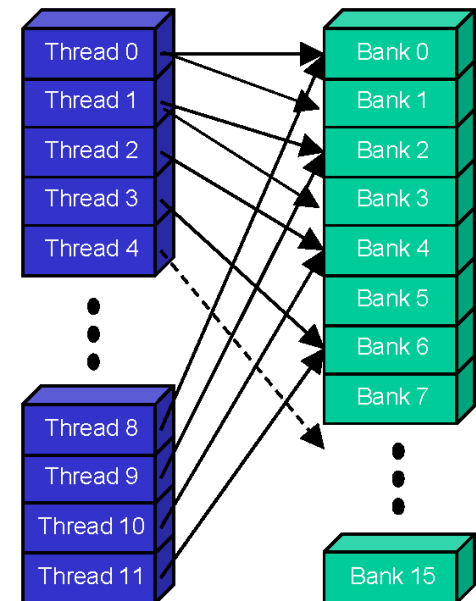
Common Array Bank Conflict Patterns

1D

- **Each thread loads 2 elements into shared mem:**
 - 2-way-interleaved loads result in 2-way bank conflicts:

```
int tid = threadIdx.x;  
shared[2*tid] = global[2*tid];  
shared[2*tid+1] = global[2*tid+1];
```

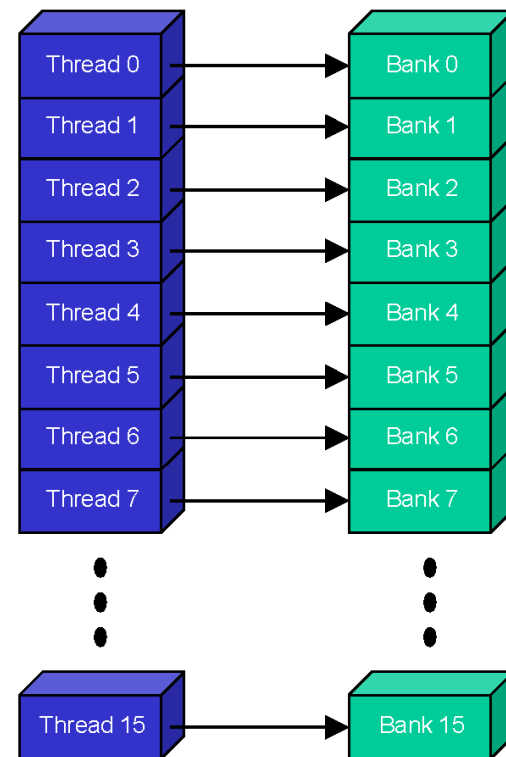
- **This makes sense for traditional CPU threads, locality in cache line usage and reduced sharing traffic.**
 - Not in shared memory usage where there is no cache line effects but banking effects



A Better Array Access Pattern

- Each thread loads one element in every consecutive group of `blockDim` elements.

```
shared[tid] = global[tid];  
shared[tid + blockDim.x] =  
    global[tid + blockDim.x];
```

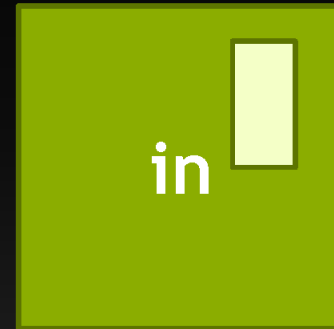


OPTIMIZE

Kernel Optimizations: *Shared Memory Accesses*

Case Study: Matrix Transpose

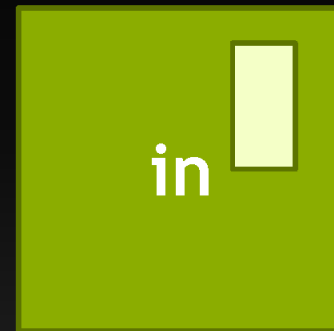
- Coalesced read
 - Scattered write (stride N)
- ⇒ Process matrix tile, not single row/column, per block
- ⇒ Transpose matrix tile within block



Case Study: Matrix Transpose

- Coalesced read
- Scattered write (stride N)
- Transpose matrix tile within block

⇒ **Need threads in a block to cooperate:
use shared memory**



Transpose with coalesced read/write

```
__global__ transpose(float in[], float out[])
{
    __shared__ float tile[TILE][TILE];

    int glob_in = xIndex + (yIndex)*N;
    int glob_out = xIndex + (yIndex)*N;

    tile[threadIdx.y][threadIdx.x] = in[glob_in];

    __syncthreads();

    out[glob_out] = tile[threadIdx.x][threadIdx.y];
}
```

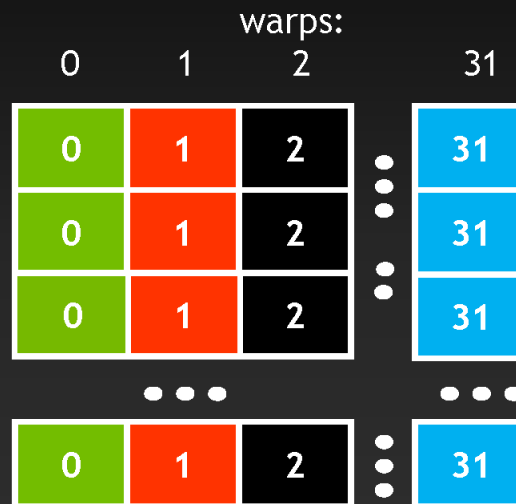
Fixed GMEM coalescing, but introduced SMEM bank conflicts

```
transpose<<<grid, threads>>>(in, out);
```

Shared Memory: Avoiding Bank Conflicts

- Example: **32x32** SMEM array
- Warp accesses a column:
 - 32-way bank conflicts (threads in a warp access the same bank)

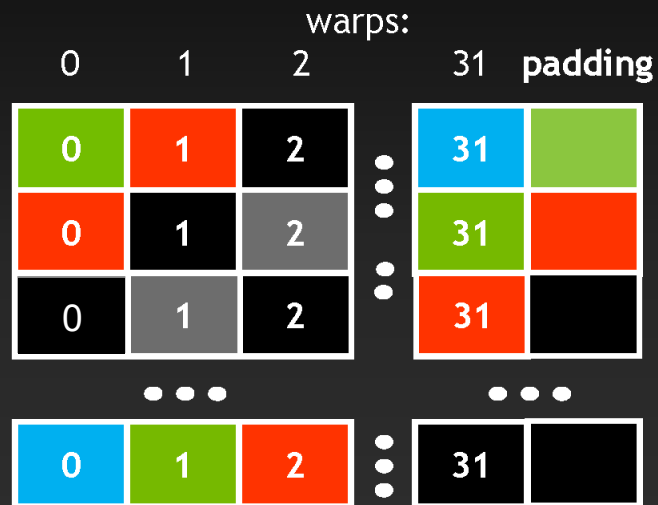
Bank 0
Bank 1
...
Bank 31



Shared Memory: Avoiding Bank Conflicts

- Add a column for padding:
 - 32x33 SMEM array
- Warp accesses a column:
 - 32 different banks, no bank conflicts

Bank 0
Bank 1
...
Bank 31



Thank you.